

Stating the Main Theorem

Yesterday, Claudia ended by sketching a functor

$$\text{Spf}: (\text{CAlg}_{\mathbb{K}}^{\text{aug}})^{\text{op}} \longrightarrow \text{Fun}(\text{CAlg}_{\mathbb{K}}^{\text{sm}}, \mathcal{S})$$

There are two versions

① Via 1-categories equipped with weak equivalences

$$\mathcal{S} = \underline{\text{sSet}} \text{ w/ wk. htpy eq [model cat.]}$$

$$\underline{\text{CAlg}}^{\text{dg}} = \text{comm dg algs over } \mathbb{K} \text{ concentrated in degrees } \leq 0 \text{ w/ quasi-iso [model cat]}$$

$$\underline{\text{CAlg}}^{\text{dg, aug}} = (\underline{\text{CAlg}}^{\text{dg}})_{/\mathbb{K}} = \left\{ \begin{array}{c} \text{obj: } A \\ \downarrow \\ \mathbb{K} \\ A \rightarrow B \\ \downarrow \circlearrowleft \\ \mathbb{K} \end{array} \right.$$

[w/ inherited model str]

$$\underline{\text{CAlg}}^{\text{small}} = \text{full subcategory of } \underline{\text{CAlg}}^{\text{dg, aug}} \text{ of small algebras [no model cat]}$$

$$\underline{\text{Fun}}(\underline{\text{CAlg}}^{\text{small}}, \underline{\text{sSet}}) = \text{functors where } X \xrightarrow{f} Y \text{ if } X(A) \xrightarrow{f(A)} Y(A) \quad \forall A \rightarrow \mathbb{K}$$

[no model cat]

Examples are often "presented" in this setting

Not sure how to distinguish substantially in a nice way: I will underline for today

② Via ∞ -categories (i.e. nice version like quasicategories)

$S' = (\infty, 1)$ -category equivalent (as $(\infty, 1)$ -cats!)
to "sSet" or "Top"

$\mathcal{C}Alg = (\infty, 1)$ -category presented by $\mathcal{C}Alg^{ob}$

(e.g. $N(\mathcal{C}Alg^{ob})$)

subset of fibrant & cofibrant objects

$\mathcal{C}Alg^{aug} =$ slice $(\infty, 1)$ -category

$\mathcal{C}Alg^{sm} =$ full subcategory of small algebras

$\text{Fun}(\mathcal{C}Alg^{sm}, S') = (\infty, 1)$ -category of functors

Theorems are easy to prove here.

Remark: One role for the abstract machinery is to relate these two settings.

In ②:

~~SP(A, R)~~

$$\text{SP} \left(\begin{array}{c} A \\ \varepsilon_A \downarrow \\ K \end{array} \right) \left(\begin{array}{c} R \\ \varepsilon_R \downarrow \\ K \end{array} \right) = \text{Map}_{\mathcal{C}Alg^{aug}} \left(\begin{array}{c} A \\ \varepsilon_A \downarrow \\ K \end{array}, \begin{array}{c} R \\ \varepsilon_R \downarrow \\ K \end{array} \right)$$

pullback in S' of one diagram
[need to talk about pullbacks!]

$$\begin{array}{c} \text{Map}_{\mathcal{C}Alg} (A, R) \\ \downarrow \varepsilon_R \\ \text{Map}_{\mathcal{C}Alg} (A, K) \end{array}$$

$\xrightarrow{\varepsilon_A}$

Let me mention that you can compute this in 1-categorical setting:

$$\text{Map}_{\text{CALg}}(A, R) \simeq \left([n] \longmapsto \text{CALg} \left(A, R \otimes \underbrace{\Omega_{\text{poly}}^*(\Delta^n)}_{\substack{\text{simplicial set} \\ \text{polynomial de Rham forms} \\ \text{on } n\text{-simplex } \Delta^n}} \right) \right)$$

$$\mathbb{K}[t_0, \dots, t_n, dt_0, \dots, dt_n] / \left(\begin{array}{l} \sum t_i = 1, \\ \sum dt_i = 0 \end{array} \right)$$

$\underbrace{\quad}_0 \quad \underbrace{\quad}_1$
 $\forall \partial(t_i) = dt_i$

"polynomial de Rham forms on n -simplex Δ^n "

As usual, you should choose A, R sufficiently nice (e.g. cofibrant & fibrant, respectively)

In a moment, we will discuss pullbacks, but if you know homotopy pullbacks of ^{top}spaces, you can mimic that here.

We also have a functor

$$C_{\text{Lie}}^* : \text{Lie}_{\mathbb{K}}^{\text{dg}} \longrightarrow (\text{CALg}_{\mathbb{K}}^{\text{dg}, \text{aug}})^{\text{op}}$$

$$(g, \partial) \longmapsto C_{\text{Lie}}^*(g) = \left(\text{Sym}_{\mathbb{K}}(g^*[-1]), d_{c^* + \partial} \right)$$

that preserves weak equivalences

$$\Rightarrow C_{\text{Lie}}^* : \text{Lie}_{\mathbb{K}} \longrightarrow (\text{CALg}_{\mathbb{K}}^{\text{aug}})^{\text{op}}$$

a functor of $(\infty, 1)$ -categories

augmentation is projection onto $\text{Sym}^0 \cong \mathbb{K}$

We can now state the main theorem.

Theorem

The composite functor

$$\Psi = \text{Spf} \circ C_{\text{Lie}}^* : \text{Lie}_{\mathbb{K}} \longrightarrow \text{Fun}(\text{CAg}_{\mathbb{K}}^{\text{sm}}, \mathcal{S})$$

factors through $\text{FMP} \subseteq \quad "$

and is an equivalence of ∞ -categories

$$\text{Lie}_{\mathbb{K}} \xrightarrow{\Psi} \text{FMP}$$



recall: $\text{ho}(\Psi)$ is an equivalence

$$\& \quad \mathbb{F}_{g,h} : \text{Map}_{\text{Lie}}(g,h) \xrightarrow{\sim} \text{Map}_{\text{FMP}}(\mathbb{F}_g, \mathbb{F}_h)$$

~~is an equivalence of simplicial sets $\forall g, h$~~

To understand this functor, we certainly need to understand homotopy pullbacks better! That will be a goal of the rest of today.

We will use our improved understanding to construct the "tangent functor"

$$\Pi : \text{FMP} \longrightarrow \text{Mod}_{\mathbb{K}} \quad (\text{coming from } \text{dgVect}_{\mathbb{K}})$$

which computes the tangent complex of a moduli functor & generalizes the usual notion of a Zariski tangent space. It will "linearize" our situation & be the first step towards an inverse functor.

Pullbacks from the perspective of ∞ -categories

There are different perspectives useful for different situations

- descriptions by universal property: good for proving structural theorems! *eg. in quasicategories*
- descriptions by construction: good for building stuff! *eg. in model categories*

v1
Here is a clean abstract definition.

Let \mathcal{C}, \mathcal{D} be quasicategories

Let $F: \mathcal{D} \rightarrow \mathcal{C}$ be a functor, i.e., map of simplicial sets.

Let $\mathcal{D}^\triangleleft$ be \mathcal{D} with an "initial object" attached,

$$\text{i.e., } \Delta^0 \star \mathcal{D} \quad \left((\Delta^0 \star \mathcal{D})_n = * \cup \mathcal{D}_n \cup \bigcup_{i+j=n-1} \mathcal{D}_j \right)$$

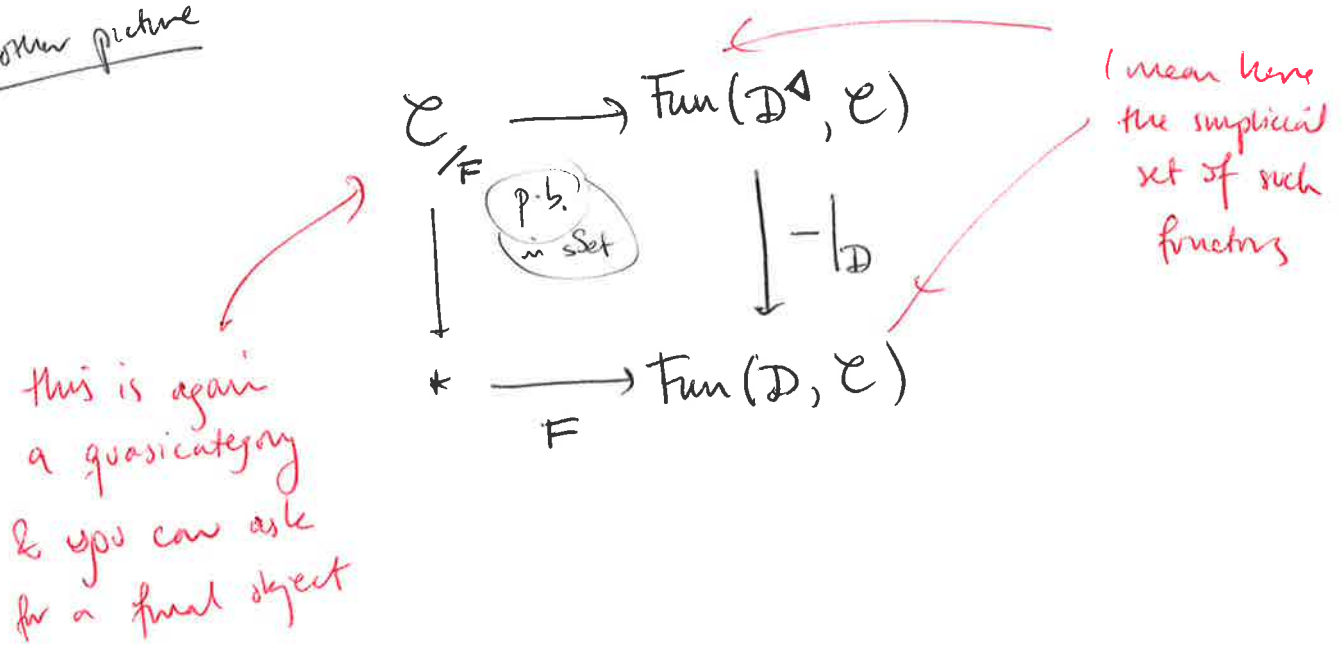
Def A limit of F is a functor $\overline{F}: \mathcal{D}^\triangleleft \rightarrow \mathcal{C}$

such that ① $\overline{F}|_{\mathcal{D}} = F$ ~~and~~

② \overline{F} is final among such extensions.

for any extension G of F to $\mathcal{D}^\triangleleft$, the space of maps $\{G \rightarrow \overline{F}\}$ is contractible

Answer picture



This notion certainly captures the idea of a limit, and the use of $s\text{Set}$ keeps track of homotopy coherence automatically.

But it may not be easy to exhibit a particular limit in practice.

v2 Suppose we have a 1-category \mathcal{C} with a subcategory of weak equivalences W .

Let \mathcal{D} be another category

then $\text{Fun}(\mathcal{D}, \mathcal{C}) = \left\{ \text{1-category of functors } \mathcal{D}^{\times} \rightarrow \mathcal{C} \right\}$

inherits a natural notion of weak equivalence:

$X \xrightarrow[\cong]{f} Y$ means $\forall d \in \mathcal{D}, f(d): X(d) \xrightarrow{\cong} Y(d)$ in \mathcal{C} .
ie $\mathcal{A} \cong \mathcal{W}$

Ex: $0 \rightarrow \begin{matrix} A \\ \downarrow \varphi \\ B \end{matrix}$ a diagram in dgVect , so φ is a chain map

This diagram is weakly equivalent to

$\text{Cone}(\text{id}_B) \xrightarrow{\varphi} B$ *could also add an acyclic complex here*
 \uparrow Their pushouts are not the same!

fix indexing

$$\begin{array}{ccc} \text{Cone}(\text{id}_B)_n = B_n \oplus B_{n+1} & \xrightarrow{\quad} & B_n \\ (b, b') & \xrightarrow{\quad} & b \\ d_{\text{cone}}(b, b') = (d_B b, b + d_B b') & & \end{array}$$

$$H_*^{\text{cone}}(\text{id}_B) = 0$$

since $d_{\text{cone}}(b, b') = 0 \Leftrightarrow b$ closed & exact w/ $d_B b' = -b$

& then $d_{\text{cone}}(b', 0) = (d_B b', b')$

Ex $* \leftarrow S^{n-1} \rightarrow *$ & $D^n \leftarrow S^{n-1} \rightarrow D^n$ are weakly equivalent
 We thus have a functor *but these pushouts are not the same:*
 $* \not\cong S^n!$

$$\begin{array}{ccc} \text{Ho}(\text{Fun}(\mathcal{D}, \mathcal{C})) & \xleftarrow{\Delta} & \text{Ho}(\mathcal{C}) \\ \downarrow & & \downarrow \\ (d \xrightarrow{\Delta_c} c) & \xleftarrow{\quad} & [c] \\ \forall d \in \mathcal{D} & & \end{array}$$

Def A homotopy limit is a right adjoint of Δ :
 $\text{Ho-lim}_{\mathcal{D}} : \text{Ho}(\text{Fun}(\mathcal{D}, \mathcal{C})) \longrightarrow \text{Ho}(\mathcal{C})$

In other words,

$$\text{Ho}(\mathcal{C}) \left([c], \text{Ho-lim}([F]) \right) \cong \text{Ho}(\text{Fun}(\mathcal{D}, \mathcal{C})) \left([\Delta], [F] \right)$$

this is the "cone definition" again but for all
D-diagrams at the same time "global"

When \mathcal{C} is a model category, one can construct

a representative of $\text{Ho-lim}(F)$ in the

functor category itself (fine print: there are

some technical conditions on \mathcal{C} and/or \mathcal{D}

to make sure you can do this but
for pullbacks, things are ok)

Ex $\mathcal{D} = \{ a \rightarrow b \leftarrow c \}$ so

$$\text{Fun}(\mathcal{D}, \mathcal{C}) = \left\{ \begin{array}{c} X(c) \\ \downarrow \\ X(a) \rightarrow X(b) \end{array} \right\}$$

Prop This functor category admits a model
structure where $f: X \rightarrow Y$ is

- a weak equivalence if f is an objectwise
weak equivalence in \mathcal{C}

$$f(a): X(a) \xrightarrow{\cong} Y(a) \dots \\ \text{in } \mathcal{C}$$

• a cofibration if f is an objectwise cofibration

• a fibration if

1) $f(b): X(b) \rightarrow Y(b)$ is a fibration

2) the canonical map

$$\begin{array}{ccc}
 X(a) & & \\
 \searrow & & \\
 Y(a) \times_{Y(b)} X(b) & \rightarrow & X(a) \\
 \downarrow & & \downarrow \\
 Y(a) & \rightarrow & Y(b)
 \end{array}$$

is a fibration

3) likewise $X(c) \rightarrow Y(c) \times_{Y(b)} X(b)$ is a fibration

Argument is straightforward: restriction to b is just C & then one "extends" to a & c . Try (e.g.) lifting properties.

Prop The functor

$$\Delta: C \rightarrow \text{Fun}(D, C)$$

$$x \mapsto (a, b, c \xrightarrow{\Delta_x} x)$$

preserves cofibrations & acyclic cofibrations (and vice versa)

~~so its left adjoint Lim_D is Quillen.~~

so its right adjoint Lim_D is Quillen.

Def The homotopy pullback of $F \in \text{Fun}(D, C)$

is $\mathbb{R}\text{Lim}_D(F)$, so can be computed by replacing F by a fibrant diagram

Extended example

Let's prove, as a key example, that $\mathbb{K}[x]/(x^3)$ is a homotopy pullback of the diagram

$$\begin{array}{ccc} & \mathbb{K} & \\ & \downarrow g & \\ \mathbb{K}[x]/(x^2) & \xrightarrow{f} & \mathbb{K}[\varepsilon_1] \end{array} \quad \textcircled{\star}$$

deg 0
↓
deg -1

This is a model example of an elementary extension.

Basic procedure:

- find a "fibrant replacement" $\widetilde{\textcircled{\star}}$
the diagram $\textcircled{\star}$
- compute the ordinary ("naive") pullback of the diagram $\widetilde{\textcircled{\star}}$

(i) Recall in $\text{CAlg}_{\mathbb{K}}^{\text{dg}}$, a map $A \xrightarrow{\varphi} B$ of cdgas is

- a weak eq if φ is quasi-isomorphism
- a fibration if φ is degreewise surjective

Every cda is thus fibrant, so $\mathbb{K}[\varepsilon]_{-1}$ is fibrant

(ii) Let's replace f by a fibration:

$$R := (\mathbb{K}[x, \delta], \partial: \delta \mapsto x^2)$$

$$\partial(x^n \delta) = x^{n+2} \Rightarrow \ker \partial = 0$$

$$\& \text{coker } \partial = \mathbb{K}[x]/(x^2)$$

$$\Rightarrow R \simeq \mathbb{K}[x]/(x^2)$$

$$\begin{array}{ccc} \tilde{f} \downarrow & \begin{array}{c} x \\ \downarrow \\ 0 \end{array} & \begin{array}{c} \delta \\ \downarrow \\ \varepsilon_{-1} \end{array} \\ \mathbb{K}[\varepsilon_{-1}] & & \end{array}$$

By construction, \tilde{f} is a levelwise surjection

(iii) let's replace g by a fibration:

$$R' = (\mathbb{K}[s, y], \partial'(s) = y) \& y^2 = 0 = y \delta$$

$$\begin{array}{ccc} \tilde{g} \downarrow & \begin{array}{c} y \\ \downarrow \\ 0 \end{array} & \begin{array}{c} \delta \\ \downarrow \\ \varepsilon_{-1} \end{array} \\ \mathbb{K}[\varepsilon_{-1}] & & \end{array}$$

(iv) What is the ordinary pullback

$$S = R \times_{\mathbb{K}[z_{-1}]} R' \quad ?$$

In degree 0, we have

$$\begin{array}{ccc} & \mathbb{K} \oplus \mathbb{K}y & y \\ & \downarrow & \downarrow \\ \mathbb{K}[x] & \longrightarrow & \mathbb{K} \\ x & \longmapsto & 0 \end{array}$$

so $S^0 = \mathbb{K}[x] \oplus \mathbb{K}y$ where $x \cdot y = 0$

$$\left(\mathbb{K}[x] \times_{\mathbb{K}} (\mathbb{K} \oplus \mathbb{K}y) \right) \cong \mathbb{K}[x, y] / (y^2, xy)$$

In degree 1, we have

$$\begin{array}{ccc} & \mathbb{K}s & \\ & \downarrow & \\ \mathbb{K}[x]s & \longrightarrow & \mathbb{K}z_{-1} \end{array}$$

so $S^1 = \mathbb{K}[x]s$

$$\begin{array}{l} \longrightarrow : x^n s \longmapsto \begin{cases} s & \text{if } n=0 \\ 0 & \text{else} \end{cases} \\ \downarrow : x^n s \longmapsto x^n s \end{array}$$

And S has differential

$$\begin{aligned} d(\cancel{s}) &= \partial(\cancel{s}) + \partial'(\cancel{s}) \\ &= x^2 + y \end{aligned}$$

$$n \geq 1: d(x^n s) = \partial(x^n s) + \partial'(x^n s) = x^{n+2}$$

hence $\text{coker } d \cong \mathbb{K}[x, y] / (x^3, xy, x^2 + y^2)$

$$\cong \mathbb{K}[x] / (x^3)$$

Discuss the generalizations: $\mathbb{K}[x]/(x^n)$ & elementary extensions

Concluding remarks

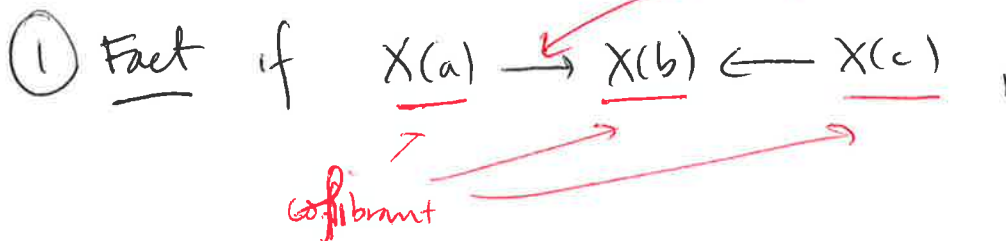


Computing homotopy pullbacks can be hard work!

Computing homotopy limits for complicated diagrams is usually awful!

(But this can be true of ordinary limits)

Some tricks can help



then naive pullback is a homotopy pullback

You don't need to replace both maps by fibrations

This fact implies

aka homotopy kernel of ϕ
or " " " " fiber " " "

Cor For $\phi: A \rightarrow B$ in dgVect , the homotopy pullback of $0 \rightarrow B \xleftarrow{\phi} A$ is presented by

$$\begin{array}{ccc} \text{Cone}(\phi)^n = (A^n \oplus B^{n+1}) & & (a, b) \\ \partial \downarrow & & \downarrow \\ \text{Cone}(\phi)^{n+1} & & (d_n a, \phi(a) + (-1)^n d_B b) \\ \downarrow & & \end{array}$$

Note. $\text{Cone}(\phi) \rightarrow A$
 $(a, b) \mapsto a$

and there is a classical homological algebra coincides \rightarrow "homotopical" algebra!

LES in $\dots \rightarrow H^n \text{Cone}(\phi) \rightarrow H^n A \rightarrow H^n B \rightarrow H^{n+1} \text{Cone}(\phi) \rightarrow \dots$

Pf replace 0 by $\text{Cone}(\text{id}_B) \simeq 0$ & take naive pullback as $\text{Cone}(\text{id}_B) \rightarrow B$ is fibration \square

② The forgetful functor $\text{CAlg}_{\mathbb{K}}^{\text{dg}} \xrightarrow{\text{fgt}} \text{dgVect}_{\mathbb{K}}$ preserves limits & is a right Quillen functor so you can compute homotopy limits in dgVect , which is often "more obvious"